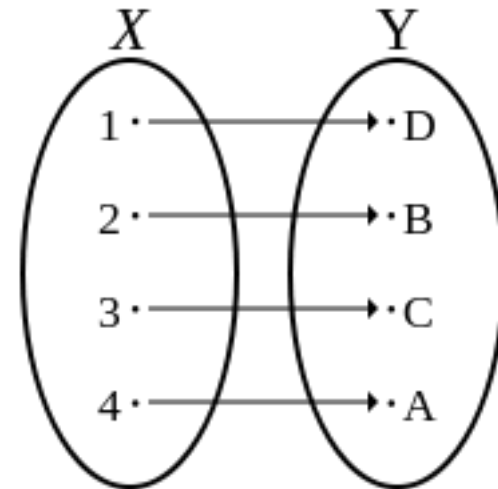
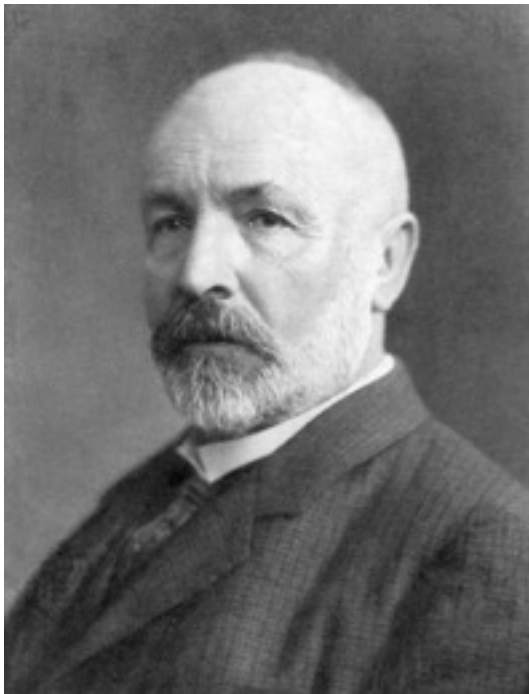


15-251

Great Theoretical Ideas in Computer Science

Lecture 5: Cantor's Legacy



September 15th, 2015

Poll

Select the ones that apply to you:

- I know what an uncountable set means.
- I know Cantor's diagonalization argument.
- I used to know what uncountable meant, I forgot.
- I used to know the diagonalization argument, I forgot.
- I've never learned about uncountable sets.
- I've never learned about the diagonalization argument.

This Week

All languages

Decidable languages

?

Factoring

$0^n | n$

Regular languages

Primality

EvenLength

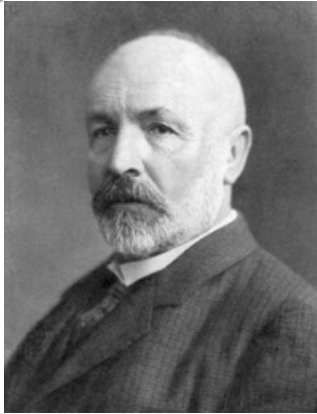
⋮

⋮

Our heroes for this week

father of set theory

father of computer science



1845-1918



1912-1954

and beyond

Uncountability

Uncomputability

Infinity in mathematics

Pre-Cantor:

“Infinity is nothing more than a figure of speech which helps us talk about limits.

The notion of a **completed infinity** doesn't belong in mathematics”

- *Carl Friedrich Gauss*



Post-Cantor:

Infinite sets are mathematical objects just like finite sets.

Some of Cantor's contributions

- > The study of infinite sets
- > Explicit definition and use of 1-to-1 correspondence
 - This is the right way to compare the cardinality of sets
- > There are different levels of infinity.
 - There are infinitely many infinities.
- > $|\mathbb{N}| < |\mathbb{R}|$ even though they are both infinite.
- > $|\mathbb{N}| = |\mathbb{Z}|$ even though $\mathbb{N} \subsetneq \mathbb{Z}$.
- > The diagonal argument.

Reaction to Cantor's ideas at the time

Most of the ideas of Cantorian set theory
should be banished from mathematics
once and for all!

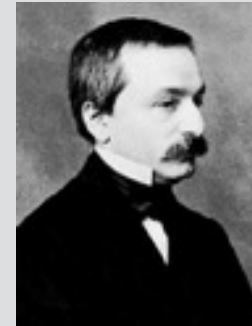
- *Henri Poincaré*



Reaction to Cantor's ideas at the time

I don't know what predominates
in Cantor's theory -
philosophy or theology.

- *Leopold Kronecker*



Reaction to Cantor's ideas at the time

Scientific charlatan.

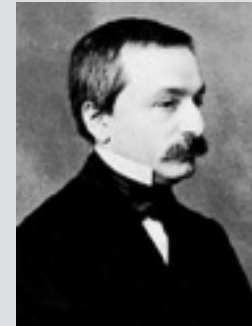
- *Leopold Kronecker*



Reaction to Cantor's ideas at the time

Corrupter of youth.

- *Leopold Kronecker*



Reaction to Cantor's ideas at the time

Wrong.

- *Ludwig Wittgenstein*



Reaction to Cantor's ideas at the time

Utter non-sense.

- *Ludwig Wittgenstein*



Reaction to Cantor's ideas at the time

Laughable.

- *Ludwig Wittgenstein*



Reaction to Cantor's ideas at the time

No one should expel us from the Paradise
that Cantor has created.

- *David Hilbert*



Reaction to Cantor's ideas at the time

If one person can see it as a paradise,
why should not another see it as a joke?

- *Ludwig Wittgenstein*



First we start with finite sets

How do we count a finite set?

$A = \{\text{apple, orange, banana, melon}\}$

What does $|A| = 4$ mean?

There is a **1-to-1 correspondence (bijection)** between

A and $\{1, 2, 3, 4\}$

apple \longleftrightarrow 1

orange \longleftrightarrow 2

banana \longleftrightarrow 3

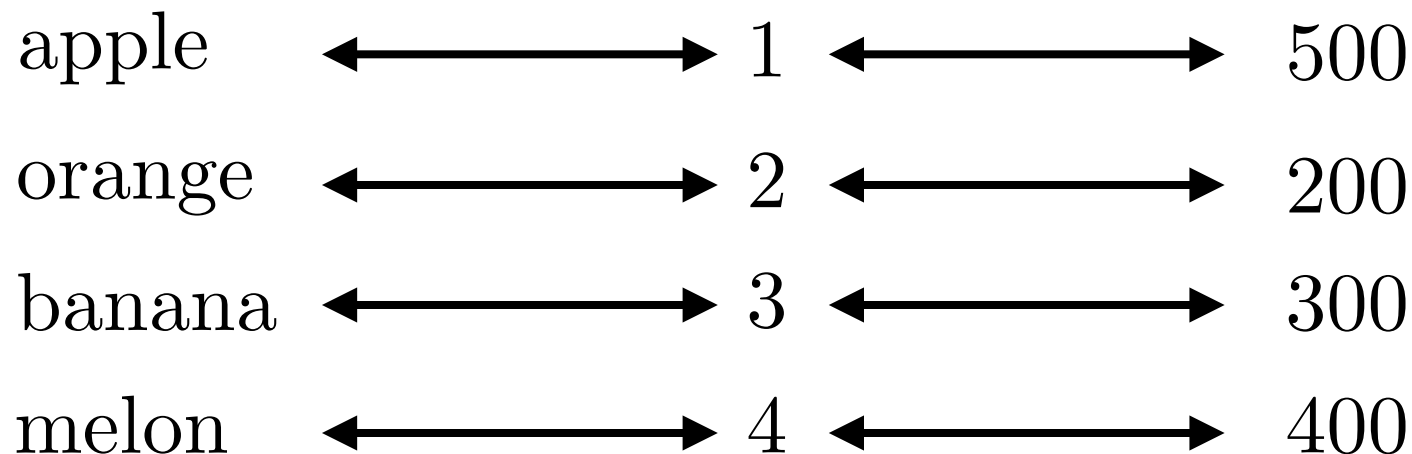
melon \longleftrightarrow 4

How do we compare the sizes of finite sets?

$A = \{\text{apple, orange, banana, melon}\}$

$B = \{200, 300, 400, 500\}$

What does $|A| = |B|$ mean?



How do we compare the sizes of finite sets?

$A = \{\text{apple, orange, banana, melon}\}$

$B = \{200, 300, 400, 500\}$

What does $|A| = |B|$ mean?

apple \longleftrightarrow 500

orange \longleftrightarrow 200

banana \longleftrightarrow 300

melon \longleftrightarrow 400

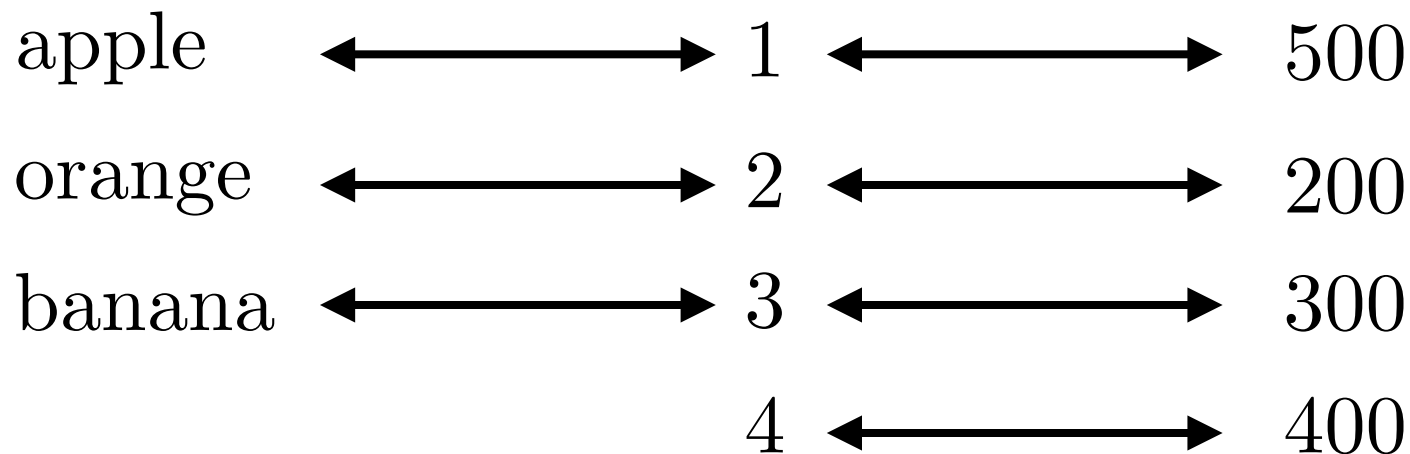
$|A| = |B|$ iff there is a **1-to-1 correspondence (bijection)** between A and B .

How do we compare the sizes of finite sets?

$A = \{\text{apple, orange, banana}\}$

$B = \{200, 300, 400, 500\}$

What does $|A| \leq |B|$ mean?



How do we compare the sizes of finite sets?

$A = \{\text{apple, orange, banana}\}$

$B = \{200, 300, 400, 500\}$

What does $|A| \leq |B|$ mean?

apple \longrightarrow 500

orange \longrightarrow 200

banana \longrightarrow 300

400

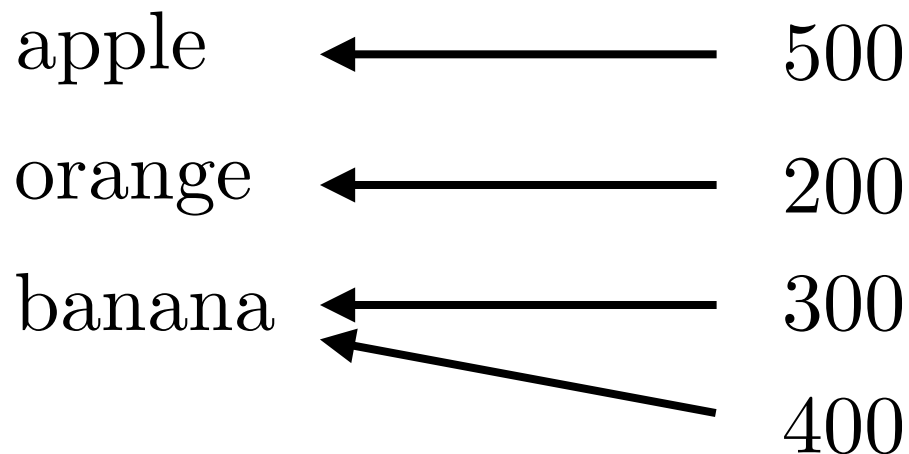
$|A| \leq |B|$ iff there is an **injection**
from A to B .

How do we compare the sizes of finite sets?

$A = \{\text{apple, orange, banana}\}$

$B = \{200, 300, 400, 500\}$

What does $|A| \leq |B|$ mean?



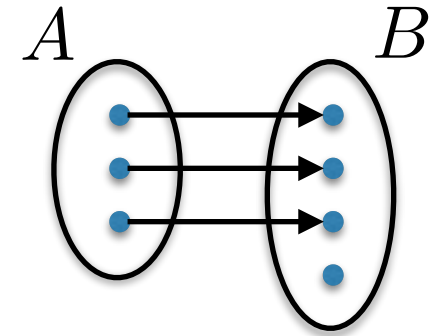
$|A| \leq |B|$ iff there is a **surjection**
from B to A .

3 important types of functions

injective, 1-to-1

$f : A \rightarrow B$ is injective if
 $a \neq a' \implies f(a) \neq f(a')$

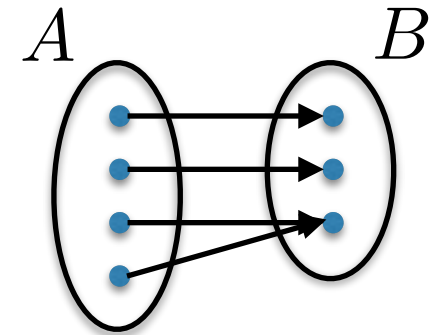
$$A \hookrightarrow B$$



surjective, onto

$f : A \rightarrow B$ is surjective if
 $\forall b \in B, \exists a \in A$ s.t. $f(a) = b$

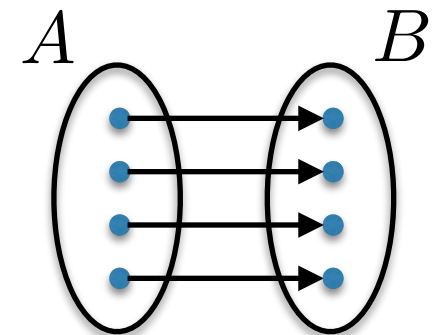
$$A \twoheadrightarrow B$$



bijective, 1-to-1 correspondence

$f : A \rightarrow B$ is bijective if
 f is injective and surjective

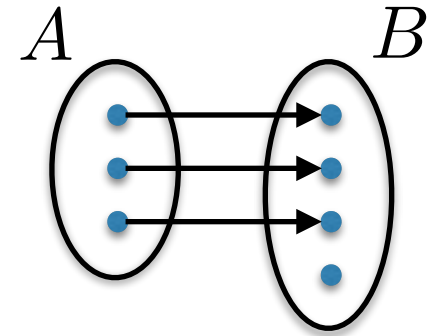
$$A \leftrightarrow B$$



Comparing the cardinality of finite sets

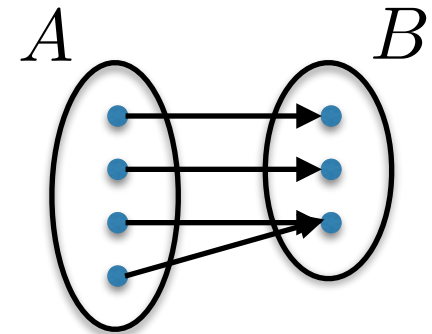
$$|A| \leq |B|$$

$$A \hookrightarrow B$$



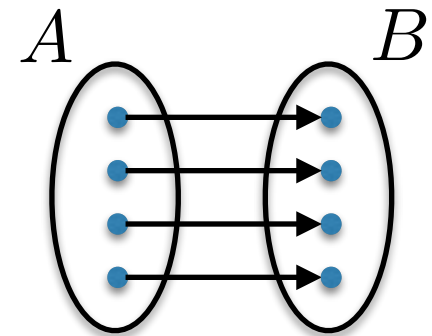
$$|A| \geq |B|$$

$$A \twoheadrightarrow B$$



$$|A| = |B|$$

$$A \leftrightarrow B$$



Sanity checks

$$|A| \leq |B| \text{ iff } |B| \geq |A|$$

$$A \hookrightarrow B \text{ iff } B \twoheadrightarrow A$$

$$|A| = |B| \text{ iff } |A| \leq |B| \text{ and } |A| \geq |B|$$

$$A \leftrightarrow B \text{ iff } A \hookrightarrow B \text{ and } A \twoheadrightarrow B$$

$$A \leftrightarrow B \text{ iff } A \hookrightarrow B \text{ and } B \hookrightarrow A$$

$$\text{If } |A| \leq |B| \text{ and } |B| \leq |C| \text{ then } |A| \leq |C|$$

$$\text{If } A \hookrightarrow B \text{ and } B \hookrightarrow C \text{ then } A \hookrightarrow C$$

One more definition

$$|A| < |B|$$

not $|A| \geq |B|$

There is no **surjection** from A to B.

There is no **injection** from B to A.

There is an **injection** from A to B,
but there is no **bijection** between A and B.

So what is the big deal?

This way of comparing the size of sets
generalizes to infinite sets!



These are the right definitions
for **infinite** sets as well!

Comparing the cardinality of infinite sets

$$|A| \leq |B|$$

$$A \hookrightarrow B$$

$$|A| \geq |B|$$

$$A \rightarrow B$$

$$|A| = |B|$$

$$A \leftrightarrow B$$

Sanity checks for infinite sets

$$|A| \leq |B| \text{ iff } |B| \geq |A|$$

$$A \hookrightarrow B \text{ iff } B \twoheadrightarrow A$$

$$|A| = |B| \text{ iff } |A| \leq |B| \text{ and } |B| \leq |A|$$

$$A \leftrightarrow B \text{ iff } A \hookrightarrow B \text{ and } A \twoheadrightarrow B$$

$$A \leftrightarrow B \text{ iff } A \hookrightarrow B \text{ and } B \hookrightarrow A$$

Cantor
Schröder
Bernstein

$$\text{If } |A| \leq |B| \text{ and } |B| \leq |C| \text{ then } |A| \leq |C|$$

$$\text{If } A \hookrightarrow B \text{ and } B \hookrightarrow C \text{ then } A \hookrightarrow C$$

So what is the big deal?



Let me show you some
interesting consequences.

Examples of equal size sets

$$|\mathbb{N}| = |\mathbb{Z}|$$

$$\mathbb{N} = \{0, 1, 2, 3, 4, \dots\}$$

$$\mathbb{Z} = \{\dots, -4, -3, -2, -1, 0, 1, 2, 3, 4, \dots\}$$

0 1 2 3 4 5 6 7 8 ...

↕ ↕ ↕ ↕ ↕ ↕ ↕ ↕ ↕

0, 1, -1, 2, -2, 3, -3, 4, -4, ...

$$f(n) = (-1)^{n+1} \left\lceil \frac{n}{2} \right\rceil$$

Examples of equal size sets

$$|\mathbb{N}| = |\mathbb{Z}|$$

Does this make any sense? $\mathbb{N} \subsetneq \mathbb{Z}$

$A \subsetneq B \implies |A| < |B|$? Shouldn't $|\mathbb{N}| < |\mathbb{Z}|$?

Does renaming the elements of a set change its size? **No.**

Let's rename the elements of \mathbb{Z} :

$\{\dots, \text{banana}, \text{apple}, \text{melon}, \text{orange}, \text{mango}, \dots\}$

Let's call this set F . How can you justify saying $|\mathbb{N}| < |F|$?

Bijection is nothing more than renaming.

Examples of equal size sets

$$|\mathbb{N}| = |S|$$

$$\mathbb{N} = \{0, 1, 2, 3, 4, \dots\}$$

$$S = \{0, 1, 4, 9, 16, \dots\}$$

$$f(n) = n^2$$

Examples of equal size sets

$$|\mathbb{N}| = |P|$$

$$\mathbb{N} = \{0, 1, 2, 3, 4, \dots\}$$

$$P = \{2, 3, 5, 7, 11, \dots\}$$

$f(n) = n$ 'th prime number.

Definition: countable and uncountable sets

Definition:

- A set A is called *countable* if $|A| \leq |\mathbb{N}|$.
- A set A is called *countably infinite* if it is infinite and countable.
- A set A is called *uncountable* if it is not countable.
(so $|A| > |\mathbb{N}|$)

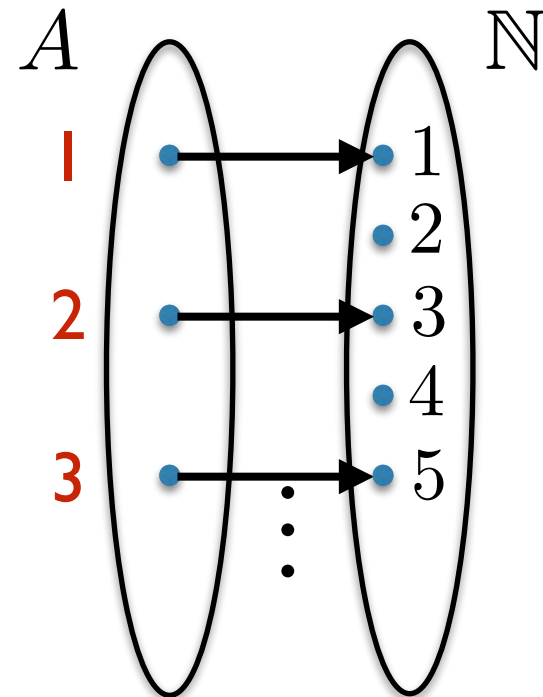
How to think about countable sets

A set A is called *countable* if $|A| \leq |\mathbb{N}|$.

So why is it called “countable”?

$|A| \leq |\mathbb{N}|$ means there is an injection $f : A \rightarrow \mathbb{N}$.

you could “count” the
elements of A
(but could go on forever)



How to think about countable sets

A set A is called *countable* if $|A| \leq |\mathbb{N}|$.

Perhaps a better name would have been *listable*:

can list the elements of A so that every element appears in the list eventually.

a_1 a_2 a_3 a_4 a_5 \dots

(this is equivalent to being countable)

How to think about countable sets

A set A is called *countable* if $|A| \leq |\mathbb{N}|$.

This seems to imply that if A is infinite, then $|A| = |\mathbb{N}|$.

Is it possible that A is infinite, but $|A| < |\mathbb{N}|$?

Theorem:

A set A is countably infinite if and only if $|A| = |\mathbb{N}|$.

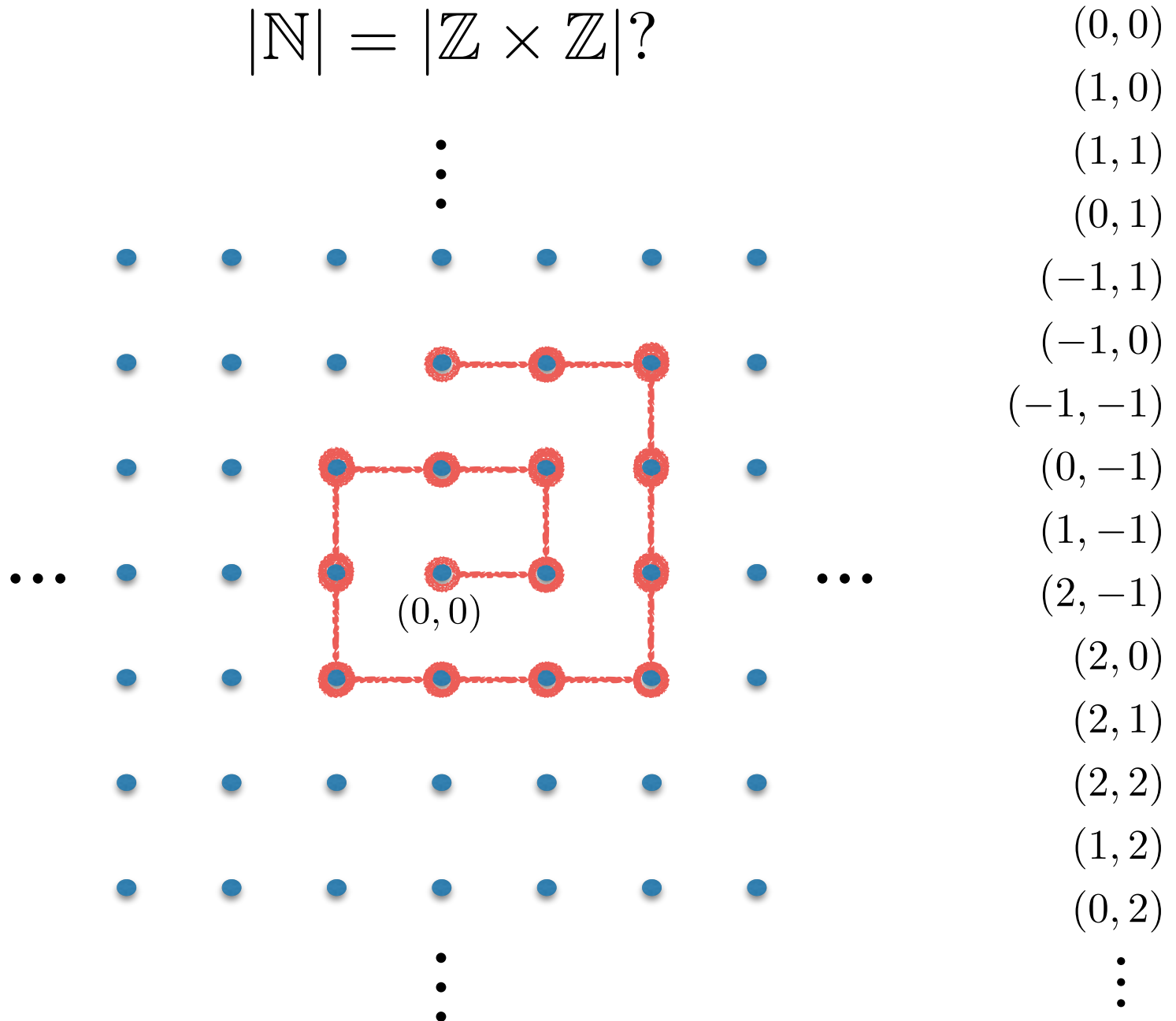
So if A is countable, there are two options:

1. A is finite
2. $|A| = |\mathbb{N}|$

Exercise: prove the theorem

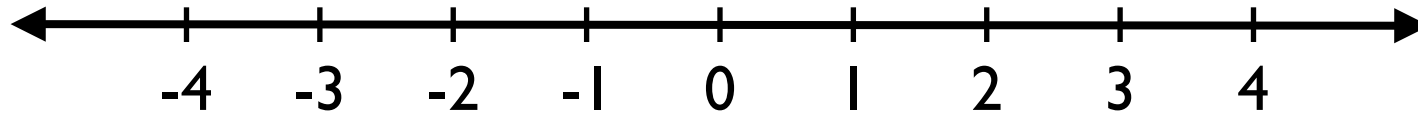
Countable?

$$|\mathbb{N}| = |\mathbb{Z} \times \mathbb{Z}|?$$



Countable?

$$|\mathbb{N}| = |\mathbb{Q}|?$$



Can we list them in the order they appear on the line?

Between **any** two rational numbers, there is another one.

Any rational number can be written as a fraction $\frac{a}{b}$.

$$\mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Q} \quad \left(\text{map } (a, b) \text{ to } \frac{a}{b} \right)$$

$$\implies |\mathbb{Q}| \leq |\mathbb{Z} \times \mathbb{Z}| = |\mathbb{N}|$$

Countable?

$$|\mathbb{N}| = |\{0, 1\}^*|?$$

$\{0, 1\}^*$ = the set of finite length binary strings.

ε

0

1

00, 01, 10, 11

000, 001, 010, 011, 100, 101, 110, 111

...

Countable?

$$|\mathbb{N}| = |\Sigma^*|?$$

Σ^* = the set of finite length words over Σ .

Same idea.

Countable?

$$|\mathbb{N}| = |\mathbb{Q}[x]|?$$

$\mathbb{Q}[x]$ = the set of polynomials with rational coefficients.

e.g. $x^3 - \frac{1}{4}x^2 + 6x - \frac{22}{7}$

Take $\Sigma = \{0, 1, \dots, 9, x, +, -, *, /, ^\}$

Every polynomial can be described by a finite string over Σ .

e.g. $x^3 - 1/4x^2 + 6x - 22/7$

So $\Sigma^* \twoheadrightarrow \mathbb{Q}[x]$ i.e. $|\mathbb{Q}[x]| \leq |\Sigma^*|$

The CS method for showing a set is countable

CS method to show a set A is countable ($|A| \leq |\mathbb{N}|$):

Show $|A| \leq |\Sigma^*|$ for some alphabet Σ .

i.e. $\Sigma^* \twoheadrightarrow A$

i.e. Show that you can encode the elements of A using finite length words over an alphabet Σ .

Seems like every set is countable...



Nope!

Cantor's Theorem

Theorem: For any non-empty set A ,

$$|A| < |\mathcal{P}(A)|.$$

$$S = \{1, 2, 3\}$$

$$\mathcal{P}(S) = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{2, 3\}, \{1, 3\}, \{1, 2, 3\}\}$$

$$|\mathcal{P}(S)| = 2^{|S|}$$

$$\mathcal{P}(S) \leftrightarrow \{0, 1\}^{|S|}$$



binary strings of length $|S|$

$$S = \{1, 2, 3\}$$

$$1 \ 0 \ 1 \ \longleftrightarrow \ \{1, 3\}$$

$$0 \ 0 \ 0 \ \longleftrightarrow \ \emptyset$$

Cantor's Theorem

Theorem: For any non-empty set A ,

$$|A| < |\mathcal{P}(A)|.$$

So:

$|\mathbb{N}| < |\mathcal{P}(\mathbb{N})|$. **I.e. $\mathcal{P}(\mathbb{N})$ is uncountable.**

$|\mathbb{N}| < |\mathcal{P}(\mathbb{N})| < |\mathcal{P}(\mathcal{P}(\mathbb{N}))| < |\mathcal{P}(\mathcal{P}(\mathcal{P}(\mathbb{N})))| < \dots$

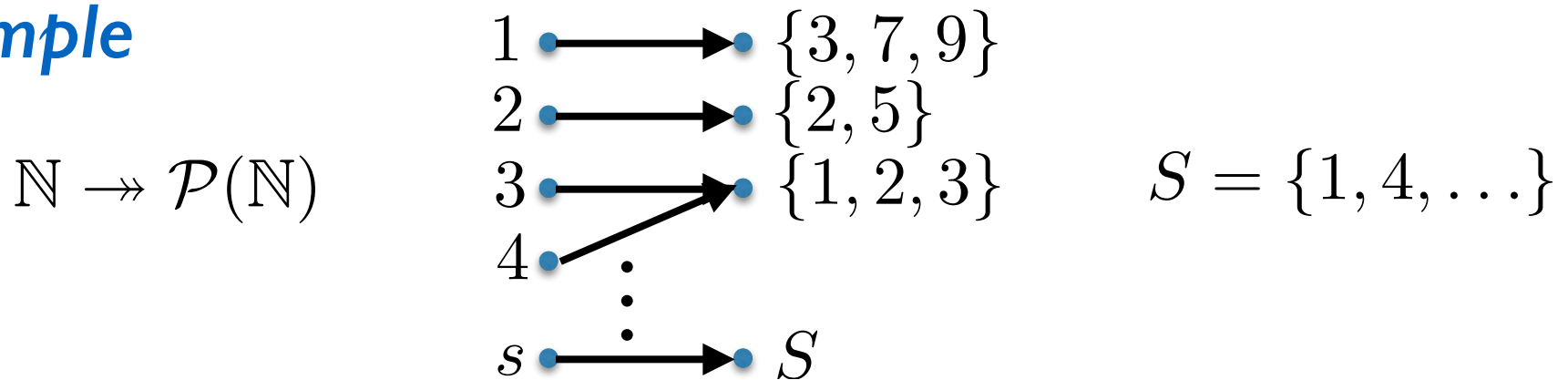
(an infinity of infinities)

Cantor's Theorem - Proof by diagonalization

Assume $|\mathcal{P}(A)| \leq |A|$ for some set A .

So $A \twoheadrightarrow \mathcal{P}(A)$. Let f be such a surjection.

Example



Define $S = \{a \in A : a \notin f(a)\} \in \mathcal{P}(A)$.

Since f is a surjection, $\exists s \in A$ s.t. $f(s) = S$.

But this leads to a contradiction:

if $s \in S$ then $s \notin f(s) = S$

if $s \notin S$ then $s \in f(s) = S$

Is $s \in S$?

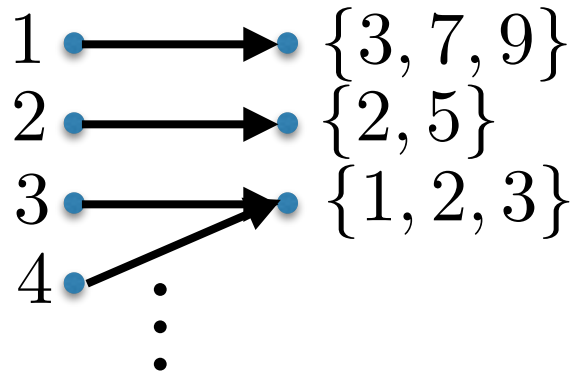
Why is this called a diagonalization argument?



Cantor's Theorem - Proof by diagonalization

Example

$$\mathbb{N} \rightarrow \mathcal{P}(\mathbb{N})$$



$$S = \{1, 4, \dots\}$$

	1	2	3	4	5	...
f(1)	0	0	1	0	0	
f(2)	0	1	0	0	1	
f(3)	1	1	1	0	0	...
f(4)	1	1	1	0	0	
f(5)	0	0	0	1	1	
\vdots			\vdots			
f(s) = S	1	0	0	1	0	...

S is defined so that S cannot equal any f(i)

Uncountable sets

Let $\{0, 1\}^\infty$ be the set of binary strings of **infinite** length.

$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \dots\}$

0000000000 ... $\longleftrightarrow \emptyset$

1111111111 ... $\longleftrightarrow \mathbb{N}$

1010101010 ... $\longleftrightarrow \{\text{even natural numbers}\}$

⋮

$\{0, 1\}^\infty$ is uncountable, i.e. $|\{0, 1\}^\infty| > |\mathbb{N}|$

because $\{0, 1\}^\infty \leftrightarrow \mathcal{P}(\mathbb{N})$. (just like $\{0, 1\}^{|S|} \leftrightarrow \mathcal{P}(S)$)

(Recall $\{0, 1\}^*$ is countable.)

Uncountable sets

Let $\{0, 1\}^\infty$ be the set of binary strings of **infinite** length.

$\{0, 1\}^\infty$ is uncountable, i.e. $|\{0, 1\}^\infty| > |\mathbb{N}|$

Direct diagonal proof: Suppose $|\{0, 1\}^\infty| \leq |\mathbb{N}|$

$\mathbb{N} \rightarrow \{0, 1\}^\infty$

1	0	0	1	0	0	...
2	0	1	0	0	1	...
3	1	1	1	0	0	...
4	1	1	1	0	0	...
5	0	0	0	1	1	...
⋮			⋮			

1 0 0 1 0 ... \rightarrow cannot appear in the list

Uncountable sets

\mathbb{R} is uncountable. In fact $(0, 1)$ is uncountable.

exercise

Be careful:

$$0.4999999999\dots = 0.5000000000\dots$$

Appreciating the diagonalization argument

If you want to appreciate something,
try to break it...



Exercise:

Why doesn't the diagonalization argument work for
 \mathbb{N} , $\{0, 1\}^*$, a countable subset of $\{0, 1\}^\infty$?

Uncountable sets

Let B be the set of bijections from \mathbb{N} to \mathbb{N} .

B is uncountable.

CS method to show a set A is uncountable ($|A| > |\mathbb{N}|$):

Show $|A| \geq |\{0, 1\}^\infty|$

i.e. $A \twoheadrightarrow \{0, 1\}^\infty$

i.e. Show that the elements of A
“encode” all the elements of $\{0, 1\}^\infty$.

One slide guide to countability questions

You are given a set A .

Is it countable or uncountable?

$$|A| \leq |\mathbb{N}| \quad \text{or} \quad |A| > |\mathbb{N}| \quad ?$$

$$|A| \leq |\mathbb{N}| :$$

- show directly that $A \hookrightarrow \mathbb{N}$ or $\mathbb{N} \twoheadrightarrow A$

- show $|A| \leq |B|$, where

$$B \in \{\mathbb{Z}, \mathbb{Z} \times \mathbb{Z}, \mathbb{Q}, \Sigma^*, \mathbb{Q}[x]\}$$

$$|A| > |\mathbb{N}| :$$

- show directly using a diagonalization argument

- show $|A| \geq |\{0, 1\}^\infty|$

An Interesting Question

Is there a set S such that

$$|\mathbb{N}| < |S| < |\mathcal{P}(\mathbb{N})|?$$

Continuum Hypothesis:

No such set exists.

(Hilbert's 1st problem)

The story continues next lecture...



and beyond