

Computational Social Choice

Anil Ada Ariel Procaccia (this time)

## SOCIAL CHOICE THEORY

- A mathematical theory that deals with aggregation of individual preferences
- Origins in ancient Greece
- Formal foundations: 18<sup>th</sup> Century (Condorcet and Borda)
- 19<sup>th</sup> Century: Charles Dodgson
- 20<sup>th</sup> Century: Nobel prizes to Arrow and Sen

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## THE VOTING MODEL

- Set of voters  $N = \{1, \dots, n\}$
- Set of alternatives A; denote |A| = m
- Each voter has a ranking over the alternatives
- Preference profile = collection of all voters' rankings

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c b

a c

c b a

а

b

## VOTE OVER CUISINES





## VOTING RULES

- Voting rule = function from preference profiles to alternatives that specifies the winner of the election
- Plurality
  - Each voter awards one point to top alternative
  - Alternative with most points wins
  - Used in almost all political elections

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### MORE VOTING RULES

#### • Borda count

- $\circ~$  Each voter awards m-k points to alternative ranked  $k'{\rm th}$
- $_{\circ}$   $\,$  Alternative with most points wins
- Proposed in the 18<sup>th</sup> Century by the chevalier de Borda
- $_{\circ}$   $\,$  Used for national elections in Slovenia
- Similar to rule used in the Eurovision song contest





Lordi Eurovision 2006 winners

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- *x* beats *y* in a pairwise election if the majority of voters prefer *x* to *y*
- Plurality with runoff
  - First round: two alternatives with highest plurality scores survive
  - Second round: pairwise election between these two alternatives

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## MORE VOTING RULES

- Single Transferable vote (STV)
  - $_{\circ}~m-1~{\rm rounds}$
  - In each round, alternative with least plurality votes is eliminated
  - $_{\circ}$   $\,$  Alternative left standing is the winner  $\,$
  - Used in:
    - Ireland, Malta, Australia, and New Zealand
    - US: Maine (governor, US congress), cities like San Francisco and Cambridge



## $\operatorname{STV}$ : EXAMPLE



### SOCIAL CHOICE AXIOMS

- How do we choose among the different voting rules? Via desirable properties!
- Majority consistency = if a majority of voters rank alternative **x** first, then **x** should be the winner
- Poll 1: Which rule is not majority consistent?
  - 1. Plurality
  - 2. Plurality with runoff
  - 3. Borda count
  - 4. STV

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## MARQUIS DE CONDORCET

- 18<sup>th</sup> Century French Mathematician, philosopher, political scientist
- One of the leaders of the French revolution
- After the revolution became a fugitive
- His cover was blown and he died mysteriously in prison



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## CONDORCET WINNER

- Recall: *x* beats *y* in a pairwise election if a majority of voters rank *x* above *y*
- Condorcet winner beats every other alternative in pairwise election

1	2	3	
а	с	b	
b	а	с	
с	b	а	

• Condorcet paradox = cycle in majority preferences



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### CONDORCET CONSISTENCY

- Condorcet consistency = select a Condorcet winner if one exists
- Poll 2: Which rule is Condorcet consistent?
  - 1. Plurality
  - 2. Borda count
  - 3. Both
  - 4. Neither

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### MORE VOTING RULES

- Copeland
  - $\circ~$  Alternative's score is #alternatives it beats in pairwise elections
  - Why does Copeland satisfy the Condorcet criterion?



## AWESOME EXAMPLE

- Plurality: *a*
- Borda: **b**

- Dorua. D	voters	voters	voters	voters	voters	voters
• Condorcet	а	b	с	с	d	е
winner: <i>c</i>	b	d	d	е	е	С
• STV: <i>d</i>	с	с	b	b	с	b
• 51 V: a	d	е	а	d	b	d
<ul> <li>Plurality</li> </ul>	е	а	е	а	а	а
with runoff:						

- Pl
  - е

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# MANIPULATION

- Using Borda count
- Top profile: **b** wins
- Bottom profile: a wins
- By changing his vote, voter 3 achieves a better outcome!
- Borda's response: "My scheme is intended only for honest men!"

d	d	d
1		3
b	b	а
а	а	с
с	С	d

d d b

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b b a

a a b

c c c

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#### **STRATEGYPROOFNESS**

- A voting rule is strategyproof (SP) if a voter can never benefit from lying about his preferences
- Poll 3: What is the largest value of m for which plurality is SP?
  - m = 1
  - 2. m = 2
  - 3. m = 3
  - $m = \infty$

## STRATEGYPROOFNESS

- A voting rule is dictatorial if there is a voter who always gets his most preferred alternative
- A voting rule is constant if the same alternative is always chosen
- Constant functions and dictatorships are SP



Constant function



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# GIBBARD-SATTERTHWAITE

- A voting rule is onto if any alternative can win
- Theorem (Gibbard-Satterthwaite): If  $m \ge 3$  then any voting rule that is SP and onto is dictatorial
- In other words, any voting rule that is onto and nondictatorial is manipulable



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## COMPLEXITY OF MANIPULATION

- Manipulation is always possible in theory
- But can we design voting rules where it is difficult in practice?
- Are there "reasonable" voting rules where manipulation is a hard computational problem? [Bartholdi et al. 1989]



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### THE COMPUTATIONAL PROBLEM

- *f*-MANIPULATION problem:
  - $\circ~$  Given votes of nonmanipulators and a preferred alternative p
  - Can manipulator cast vote that makes puniquely win under f?
- Example: Borda, p = a

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b

а

с

d

а

с

d

b

а

с

d

## A GREEDY ALGORITHM

- Rank p in first place
- While there are unranked alternatives:
  - $\circ~$  If there is an alternative that can be placed in next spot without preventing p from winning, place this alternative
  - $_{\circ}$   $\,$  Otherwise return false



	1							
1	2	3	1	2	3	1	2	3
b	b	а	b	b	а	b	b	а
а	а		а	а	b	а	а	с
с	с		с	с		С	с	
d	d		d	d		d	d	
1	2	3	1	2	3	1	2	3
b	b	а	b	b	а	b	b	а
а	а	с	а	а	с	а	а	с
с	С	b	С	с	d	С	с	d
d	d		d	d		d	d	b
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## EXAMPLE: BORDA



## WHEN DOES THE ALG WORK?

- Fact: The greedy algorithm is a polynomial-time algorithm for R-MANIPULATION for  $R \in \{\text{plurality, Borda count, plurality with runoff, Copeland,...}$
- Theorem [Bartholdi and Orlin, 1991]: The STV-MANIPULATION problem is NP-complete!



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### IS SOCIAL CHOICE PRACTICAL?

- UK referendum: Choose between plurality and STV as a method for electing MPs
- Academics agreed STV is better...
- ... but STV seen as beneficial to the hated Nick Clegg
- Hard to change political elections!

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### COMPUTATIONAL SOCIAL CHOICE

#### • However:

- in online voting...
- in human computation...
- in multiagent systems...

the designer is free to employ any voting rule!







## SUMMARY

- Terminology:
  - Plurality, Borda count, plurality with runoff, STV, Copeland
  - Majority consistency
  - Condorcet winner, Condorcet consistency
  - $\circ$  Strategyproofness
  - The Gibbard-Satterthwaite Thm
- Principles:
  - $_{\circ}$   $\,$  NP-hardness can be good!

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