

Some motivating real-world examples
matching professors and courses


15-1 10
$15-112$
15-122
15-150
|5-25|
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| Some motivating real-world examples |  |
| :---: | :---: |
| matching rooms and courses |  |
| GHC 440I | $15-110$ |
| DH 2210 | $15-112$ |
| GHC 5222 | $15-122$ |
| WEH 7500 | $15-150$ |
| DH 2315 | $15-251$ |
| $\vdots$ | $\vdots$ |



How do you solve a problem like this?
I. Formulate the problem
2. Ask: Is there a trivial algorithm?
3. Ask: Is there a better algorithm?
4. Find and analyze

## Remember the CS life lesson

## First step: Formulate the problem

## Purpose:

- Get rid of all the distractions, identify the crux.
- Get a clean mathematical model that is easier to reason about.
- Solutions often generalize to other settings.


## Bipartite Graphs


$G=(V, E)$ is bipartite if:

## Bipartite Graphs

Given a graph $G=(V, E)$, we could ask, is it bipartite?



## Important Characterization

An obstruction for being bipartite:
Contains a cycle of odd length.

Is this the only type of obstruction?

## Theorem:

## Bipartite Graphs



Often we write the bipartition explicitly:

$$
G=(X, Y, E)
$$

## Bipartite Graphs

Great at modeling relations between two classes of objects.

## Examples:

$X=$ machines, $Y=$ jobs
An edge $\{x, y\}$ means $x$ is capable of doing $y$.
$X=$ professors, $Y=$ courses
An edge $\{x, y\}$ means $x$ can teach $y$.
$X=$ students, $Y=$ internship jobs
An edge $\{x, y\}$ means $x$ and $y$ are interested in each other.
:

## Matchings in bipartite graphs

Often, we are interested in finding a matching in a bipartite graph
matching


## A matching :

## Matchings in bipartite graphs

Often, we are interested in finding a matching in a bipartite graph


## Maximum matching:

## Matchings in bipartite graphs

Often, we are interested in finding a matching in a bipartite graph
maximal matching


## Maximal matching:

## Matchings in bipartite graphs

Often, we are interested in finding a matching in a bipartite graph
perfect matching


## Perfect matching:

How many different perfect matchings does the graph have (in terms of n )?


$$
|X|=|Y|=n
$$

## Important Note

We can define matchings for non-bipartite graphs as well.


## Maximum matching problem

The problem we want to solve is:

## Maximum matching problem

Input: A graph $G=(V, E)$.
Output: A maximum matching in $G$.

## Bipartite maximum matching problem

Actually, we want to solve the following restriction:

## Bipartite maximum matching problem

Input: A bipartite graph $G=(X, Y, E)$.
Output: A maximum matching in $G$.

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## Bipartite maximum matching problem

Bipartite maximum matching problem
Input: A bipartite graph $G=(X, Y, E)$.
Output: A maximum matching in $G$.

Is there a (trivial) algorithm to solve this problem?

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Bipartite maximum matching problem
A good first attempt:
What if we picked edges "greedily"?


## Bipartite maximum matching problem

A good first attempt:
What if we picked edges "greedily"?


Is there a way to get out of this local optimum?

## Important Definition: Augmenting paths

Let $M$ be some matching.
An alternating path with respect to $\mathbf{M}$ is a path in $\mathbf{G}$ such that:


An augmenting path with respect to $M$ is an alternating path such that:

Important Definition: Augmenting paths


Important Definition: Augmenting paths
 matching $=$ red edges

Augmenting path:
2-5-I-7

augmenting path $\Longrightarrow$ can obtain a bigger matching.

## Important Definition: Augmenting paths


matching $=$ red edges
Augmenting path:
4-8

augmenting path $\Longrightarrow$ can obtain a bigger matching.

Augmenting paths and maximum matchings
augmenting path $\Longrightarrow$ can obtain a bigger matching.

## In fact, it turns out:

no augmenting path $\Longrightarrow$ maximum matching.

## Theorem:

## Augmenting paths and maximum matchings

## Proof:

If there is an augmenting path with respect to $\mathbf{M}$, we saw that $M$ is not maximum.

## Want to show:

If $M$ not maximum, there is an augmenting path w.r.t. M.
Let $M^{*}$ be a maximum matching. $|M *|>|M|$.


Let $\mathbf{S}$ be the set of edges contained in $\mathbf{M}^{*}$ or $\mathbf{M}$ but not both.
$\mathbf{S}=(\mathbf{M} * \cup \mathbf{M})-(\mathbf{M} \cap \mathbf{M} \boldsymbol{*})$

Augmenting paths and maximum matchings

## Proof (continued):



$$
\mathbf{S}=(\mathbf{M} * \cup \mathbf{M})-(\mathbf{M} \cap \mathbf{M} *)
$$


(will find an augmenting path in $\mathbf{S}$ )

Augmenting paths and maximum matchings

## Proof (continued):


$\mathbf{S}=(\mathbf{M} * \cup \mathbf{M})-(\mathbf{M} \cap \mathbf{M} *)$
(will find an augmenting path in $\mathbf{S}$ )

## Augmenting paths and maximum matchings

## Theorem:

A matching $\mathbf{M}$ is maximum if and only if there is no augmenting path with respect to $\mathbf{M}$.

## Summary of proof:

Algorithm to find maximum matching
Theorem:
A matching $M$ is maximum if and only if there is no augmenting path with respect to $\mathbf{M}$.

## Algorithm to find max matching:

Finding augmenting paths in bipartite graphs


Finding augmenting paths in bipartite graphs


## Algorithm:

## Running time:

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