Almost the Same
Recall that an independent set of vertices in a graph is a set such that no pair are adjacent. We define the INDEPENDENT-SET problem to be the problem of checking if there exists an independent set of size $k$. Prove INDEPENDENT-SET is NP-COMPLETE by reducing CLIQUE to it.

Your Covering is Satisfactory
Prove that VERTEX-COVER is NP-HARD, by reducing 3-SAT to it.

Not a Graph Problem
Prove that SUBSET-SUM, the problem of given $n$ numbers, is there some subset that sums to $k$ (we allow duplicate numbers), is NP-COMPLETE.

HALTS is HARD
Prove that the Halting Problem is NP-HARD.