

## Probability 1



## France, 1654



"Chevalier de Méré"  
AKA Antoine Gombaud

Let's bet:  
I will roll a die four times.  
I win if I get a 1.

(not actually Méré)



## France, 1654



Antoine Gombaud,  
AKA "Chevalier de Méré"

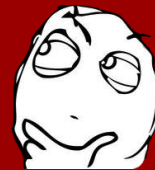
Hmm.  
No one wants to take  
this bet any more.

## France, 1654



Antoine Gombaud,  
AKA "Chevalier de Méré"

New bet:  
I will roll two dice, 24 times.  
I win if I get double-1's.



## France, 1654



Antoine Gombaud,  
AKA "Chevalier de Méré"

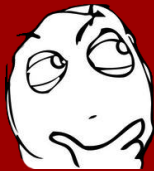
Hmm.  
I keep losing money!

## France, 1654



Alice and Bob are flipping a coin. Alice gets a point for Heads, Bob a point for Tails. First one to 4 points wins the stake of 100 francs.

Alice is leading 3-2 when gendarmes arrive to break up the game. How should they divide the stakes?



Pascal



Fermat

Probability Theory is Born

## Moral of the Story:

Analyzing gambling is not a side-benefit of probability.

Probability was invented to analyze gambling.

This is not  
"Great Theoretical Ideas  
in Gambling"

This is  
"Great Theoretical Ideas  
in Computer Science"

Probability Theory  
=  
Analyzing Code with  
Random Number Generators

“Méré throws four 6-sided dice.”



```
die1 = RandInt(6)  
die2 = RandInt(6)  
die3 = RandInt(6)  
die4 = RandInt(6)
```

“Méré throws an 8-sided die  
and a 3-sided die.”



```
die1 = RandInt(8)  
die2 = RandInt(3)
```



**def:** ‘Experiment’ = some randomized code

“A patient has a 10% chance  
of having a certain disease...”



```
x = RandInt(10)  
if x == 1 then  
    patient.hasDisease = 1  
else  
    patient.hasDisease = 0
```

“A patient has a 10% chance  
of having a certain disease...”



```
patient.hasDisease = Bernoulli(.1)
```

**Bernoulli**(p) returns 1 with probability p,  
0 with probability 1-p

“Antoine flips two fair coins.”



```
coin1 = Bernoulli(1/2)  
coin2 = Bernoulli(1/2)
```

```
{  
    if Bernoulli(1/2) == 0  
        then coin1 = Heads  
    else coin1 = Tails  
}
```

The two random generators we allow:

**RandInt** (*m*) returns 1, 2, 3, ..., *m*  
with probability  $1/m$  each

**Bernoulli** (*p*) returns 1 with probability *p*,  
0 with probability  $1-p$

**NOT ALLOWED IN 15-251:**

~~**Uniform** (*0, 1*) returns a random real number  
between 0 and 1~~

## How to Analyze Random Code

Mary flips a fair coin. If it's heads, she rolls a 3-sided die. If it's tails, she rolls a 4-sided die.

**STEP 1:** Translate to code.

```
flip = Bernoulli(1/2)
if flip == 0 (Heads) then
    die = RandInt(3)
else
    die = RandInt(4)
```

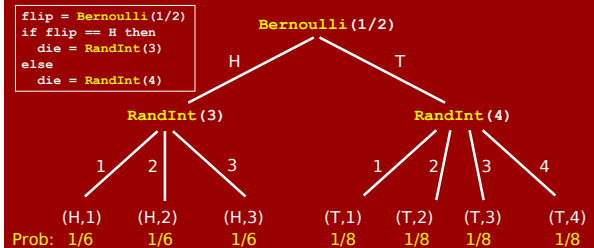
## How to Analyze Random Code

Mary flips a fair coin. If it's heads, she rolls a 3-sided die. If it's tails, she rolls a 4-sided die.

**STEP 2:** Draw a probability tree.

```
flip = Bernoulli(1/2)
if flip == 0 (Heads) then
    die = RandInt(3)
else
    die = RandInt(4)
```

## How to Analyze Random Code



Have branching for each call to a generator  
Label the leaves with **"outcomes"**  
Under each, write its probability: **multiply along the path**

### Outcome:

A leaf in the probability tree.

I.e., a possible sequence of values of all calls to generators in an execution.

### Sample Space:

The **set** of all outcomes.

E.g., { (H,1), (H,2), (H,3), (T,1), (T,2), (T,3), (T,4) }

### Probability:

Each outcome has a nonnegative probability.  
Sum of all outcomes' probabilities always 1.

## How to Analyze Random Code

Mary flips a fair coin. If it's heads, she rolls a 3-sided die. If it's tails, she rolls a 4-sided die.

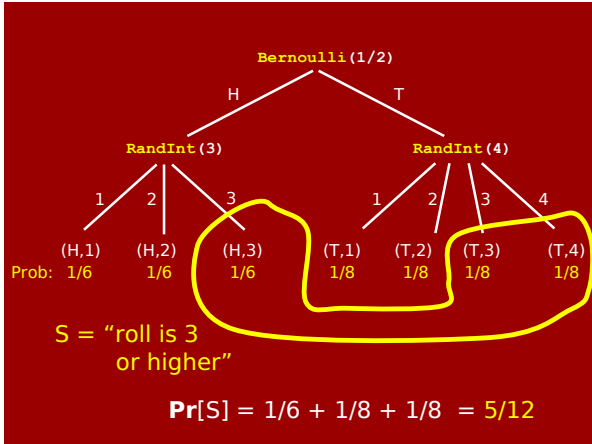
What is the probability **die roll is 3 or higher?**

### Event:

A **subset** of outcomes.

In our example,  $S = \{ (H,3), (T,3), (T,4) \}$ .

$\Pr[S]$  = sum of the probabilities of the outcomes in *S*.



## France, 1654



Alice and Bob are flipping a coin. Alice gets a point for Heads, Bob a point for Tails. First one to 4 points wins the stake of 100 francs.

Alice is leading 3-2 when gendarmes arrive to break up the game. How should they divide the stakes?

## France, 1654



It seems fair that Alice should get

(100 francs) x  $\Pr$ [Alice would win].

So let's compute that!



Alice leading 3-2:



Event A = "Alice wins" = { H, TH }

$\Pr[A] = 1/2 + 1/4 = 3/4$

## Events and Probabilities: Facts

Since  $\Pr[A]$  = sum of probs of outcomes in A, ...

If  $A \subseteq B$  then  $\Pr[A] \leq \Pr[B]$

"not A"  $\Pr[A^c] = 1 - \Pr[A]$

"A or B"  $\Pr[A \cup B] = \Pr[A] + \Pr[B] - \Pr[A \cap B]$

**FALLACY:**  $\Pr[A \cup B] = \Pr[A] + \Pr[B]$

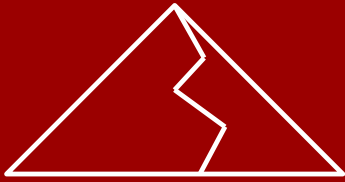
**True:**  $\Pr[A \cup B] \leq \Pr[A] + \Pr[B]$

## France, 1654



Let's bet:

I will roll a die four times.  
I win if I get a 1.



(4,6,1,2) Prob.  $1/6^4$

Let  $W$  be the event that Méré wins  
Easier to compute  $\Pr[W^c]$ .

$W^c = \{ \text{all outcomes with no 1's} \}$

$|W^c| = 5^4$

$\therefore \Pr[W^c] = 5^4/6^4$

$\therefore \Pr[W] = 1 - 5^4/6^4 \approx 51.8\%$

## France, 1654



Let's bet:  
I will roll two dice 24 times.  
I win if I get a double-1's.

$\Pr[\text{Méré wins}] =$

$$1 - 35^{24}/36^{24}$$

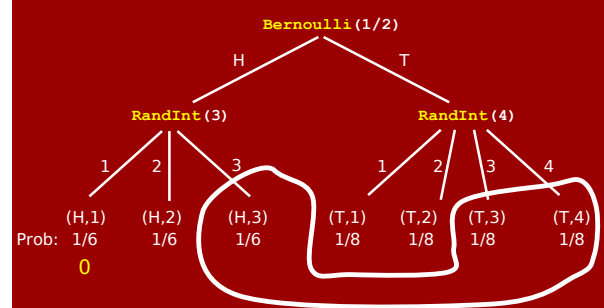
$$\approx 49.1\%$$

## Conditioning

= Revising probabilities based on 'partial information'

'Partial information' = an event

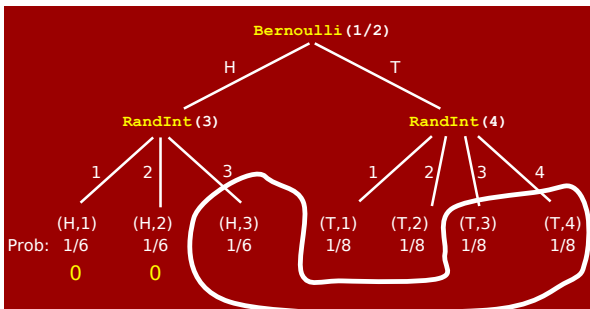
'Conditioning on event A' is like assuming/promising A occurs.



Condition on  $S$ , the event "roll is 3 or higher"

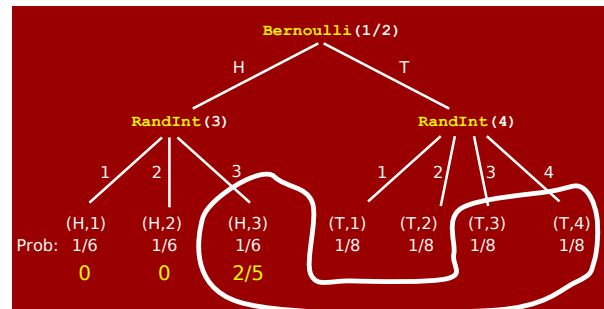
$$\Pr[(H,1) | S] = 0$$

"probability of outcome (H,1) conditioned on event S"



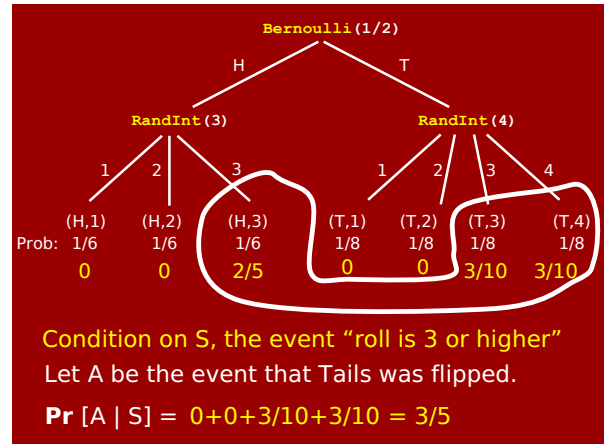
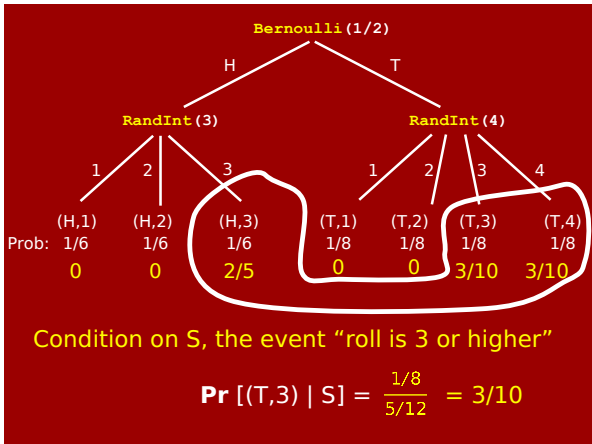
Condition on  $S$ , the event "roll is 3 or higher"

$$\Pr[(H,2) | S] = 0$$



Condition on  $S$ , the event "roll is 3 or higher"

$$\Pr[(H,3) | S] = \frac{1/6}{5/12} = 2/5$$



## Conditioning: formally

Given an experiment, let A be an event.  
(with nonzero probability)

The conditional probability of outcome  $l$  is

$$\Pr[l | A] = \begin{cases} 0 & \text{if } l \notin A, \\ \frac{\Pr[l]}{\Pr[A]} & \text{if } l \in A. \end{cases}$$

$$\therefore \Pr[B | A] = \sum_{l \in B} \Pr[l | A] = \sum_{l \in B \cap A} \frac{\Pr[l]}{\Pr[A]} = \frac{\Pr[B \cap A]}{\Pr[A]}$$

## "Chain Rule"

$$\Pr[A \cap B] = \Pr[A] \cdot \Pr[B | A]$$

"For A **and** B to occur, first A must occur (probability  $\Pr[A]$ ), and then B must occur given that A occurred (probability  $\Pr[B | A]$ )." ."

## "Chain Rule"

$$\Pr[A \cap B] = \Pr[A] \cdot \Pr[B | A]$$

$$\Pr[A \cap B \cap C] = \Pr[A] \cdot \Pr[B | A] \cdot \Pr[C | A \cap B]$$

$$\Pr[A \cap B \cap C \cap D] = \Pr[A] \cdot \Pr[B | A] \cdot \Pr[C | A \cap B] \cdot \Pr[D | A \cap B \cap C]$$

etc.

## Silver and Gold: a problem

One bag contains two silver coins.  
Another contains two gold coins.  
Another contains one silver and one gold.



Mark picks a bag at random, then picks a coin from it at random.

It turns out to be gold. What is the probability the *other* coin in his bag is gold?

## Silver and Gold: a problem



Let  $G_1$  be the event that the first chosen coin is gold.  
Let  $G_2$  be the event that the second coin in the bag is gold.  
The problem is asking us to find  $\Pr[G_2 | G_1]$ .

$$\Pr[G_1] = 3/6 = 1/2 \quad (\text{each coin equally likely to be first})$$

$$\Pr[G_1 \cap G_2] = 1/3 \quad (\text{if and only if gold-gold bag picked})$$

$$\therefore \Pr[G_2 | G_1] = \frac{\Pr[G_1 \cap G_2]}{\Pr[G_1]} = \frac{1/3}{1/2} = 2/3$$

## Law of Total Probability

or, how to actually calculate stuff

$$\Pr[B] = \Pr[A] \cdot \Pr[B | A] + \Pr[A^c] \cdot \Pr[B | A^c]$$

“Regarding event B — either A occurs  
(this has probability  $\Pr[A]$ ),  
in which case B occurs with probability  $\Pr[B | A]$ ;  
or, A does not occur  
(this has probability  $\Pr[A^c] = 1 - \Pr[A]$ ),  
in which case B occurs with probability  $\Pr[B | A^c]$ .”

## Law of Total Probability

or, how to actually calculate stuff

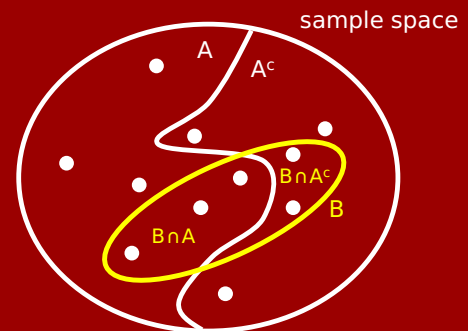
$$\Pr[B] = \Pr[A] \cdot \Pr[B | A] + \Pr[A^c] \cdot \Pr[B | A^c]$$

**Proof:**

$$\Pr[A] \cdot \Pr[B | A] = \Pr[B \cap A]$$

$$\text{Similarly, } \Pr[A^c] \cdot \Pr[B | A^c] = \Pr[B \cap A^c]$$

Each outcome in B is in exactly one of  $B \cap A$ ,  $B \cap A^c$   
Thus  $\Pr[B] = \Pr[B \cap A] + \Pr[B \cap A^c]$



## Law of Total Probability

more general version

Let events  $A_1, \dots, A_n$  be a **partition** of the sample space, meaning each outcome is in exactly one.

Then for any event B,

$$\Pr[B] = \Pr[A_1] \cdot \Pr[B | A_1] + \dots + \Pr[A_n] \cdot \Pr[B | A_n]$$

## Example

“I roll 101 regular dice. What is the probability their sum is divisible by 6?”

**Trick:** “Condition on” the sum of the first 100.





## Independence of Multiple Events

def:  $A_1, \dots, A_5$  are independent if

$$\Pr[A_1 \cap A_2 \cap A_3 \cap A_4 \cap A_5] = \Pr[A_1] \Pr[A_2] \Pr[A_3] \Pr[A_4] \Pr[A_5]$$

$$\& \Pr[A_1 \cap A_2 \cap A_3 \cap A_4] = \Pr[A_1] \Pr[A_2] \Pr[A_3] \Pr[A_4]$$

$$\& \Pr[A_1 \cap A_3 \cap A_5] = \Pr[A_1] \Pr[A_3] \Pr[A_5]$$

& in fact, the definition requires

$$\Pr \left[ \bigcap_{i \in S} A_i \right] = \prod_{i \in S} \Pr[A_i] \quad \text{for all } S \subseteq \{1, 2, 3, 4, 5\}$$

## Independence of Multiple Events

def:  $A_1, \dots, A_5$  are independent if

$$\Pr \left[ \bigcap_{i \in S} A_i \right] = \prod_{i \in S} \Pr[A_i] \quad \text{for all } S \subseteq \{1, 2, 3, 4, 5\}$$

Similar 'Principle of Independence' holds  
(5 blocks of code which don't affect each other)

Consequence: anything like

$$\Pr[A_1 \mid (A_2 \cup A_3) \cap (A_4^c \cup A_5)] = \Pr[A_1]$$

## Birthday Problem

Question:

There are  $m$  students in a room ( $m \leq 365$ ).  
What's the probability they  
all have different birthdays?

Modeling:

Ignore Feb. 29. Assume days equally likely.  
Assume no twins in the class.

```
for i = 1..m
    student[i].bday = RandInt(365)
```

## Birthday Problem — Analysis

Let  $A_i$  be event that student  $i$ 's bday differs  
from the bday of all previous students.

Let  $D$  be event that all bdays are different.

$$D = A_1 \cap A_2 \cap A_3 \cap \dots \cap A_m$$

Chain rule:

$$\Pr[D] = \Pr[A_1] \Pr[A_2|A_1] \Pr[A_3|A_1 \cap A_2] \Pr[A_4] \dots \text{etc.}]$$

So what is  $\Pr[A_i \mid A_1 \cap A_2 \cap \dots \cap A_{i-1}]$  ?

## Birthday Problem — Analysis

Let  $A_i$  be event that student  $i$ 's bday differs  
from the bday of all previous students.

So what is  $\Pr[A_i \mid A_1 \cap A_2 \cap \dots \cap A_{i-1}]$  ?

$A_1 \cap A_2 \cap \dots \cap A_{i-1}$  means first  $i-1$  students all  
had different birthdays.

$i-1$  out of 365 occupied when  $i$ th bday chosen.

$$\Pr[A_i \mid A_1 \cap A_2 \cap \dots \cap A_{i-1}] = \frac{365 - (i-1)}{365} = 1 - \frac{i-1}{365}$$

## Birthday Problem — Analysis

Let  $A_i$  be event that student  $i$ 's bday differs  
from the bday of all previous students.

Let  $D$  be event that all bdays are different.

$$\Pr[D] = \Pr[A_1] \Pr[A_2|A_1] \Pr[A_3|A_1 \cap A_2] \Pr[A_4] \dots \text{etc.}]$$

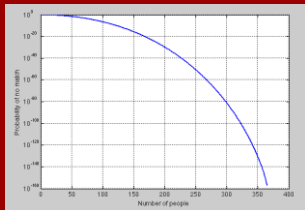
$$= 1 \cdot \left(1 - \frac{1}{365}\right) \cdot \left(1 - \frac{2}{365}\right) \dots \left(1 - \frac{m-1}{365}\right)$$

This is the final answer.

## Birthday Problem — Analysis

Pr[all m students have different bdays]

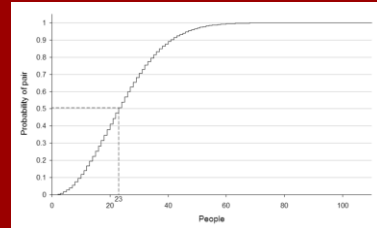
$$= 1 \cdot \left(1 - \frac{1}{365}\right) \cdot \left(1 - \frac{2}{365}\right) \cdots \left(1 - \frac{m-1}{365}\right)$$



## Birthday Problem — Analysis

Pr[in m students, some pair share a bday]

$$= 1 - 1 \cdot \left(1 - \frac{1}{365}\right) \cdot \left(1 - \frac{2}{365}\right) \cdots \left(1 - \frac{m-1}{365}\right)$$



## Birthday Problem —

Sometimes called the **Birthday “Paradox”**, because 23 seems surprisingly small.

## Birthday Problem — Analysis

What if there are N possible “birthdays”?

Pr[in m students, some pair share a “bday”]

$$= 1 - 1 \cdot \left(1 - \frac{1}{N}\right) \cdot \left(1 - \frac{2}{N}\right) \cdots \left(1 - \frac{m-1}{N}\right)$$

For what value of m is this  $\approx 1/2$  ?

This is not a calculus class, so I’ll just tell you:

$$\text{for } m \approx \sqrt{N}$$

## Birthday Problem —

Sometimes called the **Birthday “Paradox”**, because 23 seems surprisingly small.

Sometimes called the **Birthday “Attack”** in theoretical cryptography. Why...?

## Cryptographic Hash Functions

“Scrambles” any string S into a k-bit ‘hash’ f(S)

- Given f(S), should be ‘hard’ to recover S.
- Should be ‘hard’ to find a “collision”: two strings  $S_1 \neq S_2$  with  $f(S_1) = f(S_2)$ .

**Applications:** authentication schemes, data integrity schemes, digital signatures, e-cash...

## Cryptographic Hash Functions

1991: Rivest publishes MD5. (k=128)

1993: NSA publishes SHA-0. (k=160)

1995: NSA: "Um, never mind.

Please use SHA-1 instead."

SHA-1 was/is **widely used**: SSL, PGP, ...

2001: NSA also introduces SHA-2

(variants with k=224,256,384,512)

2012: Non-NSA introduces SHA-3

## Birthday Attack

Imagine trying to find a collision for SHA-1:

Take a huge number of strings, hash them all, hope that two hash to the same 160 bits.

If SHA-1 is really safe, each hash  $f(S)$  should be like  $\text{RandInt}(2^{160})$ .

This is like the Birthday Problem with  $N = 2^{160}$ !

So # tries before good chance of collision:

$$\approx \sqrt{2^{160}} = 2^{80} = 1208925819614629174706176$$

## Birthday Attack

Everybody knows this.

$2^{80}$  is considered safely "too large".

A crypto hash function is considered "broken" if you can beat the Birthday Attack.



Xiaoyun Wang (王小云)

2005: SHA-1 collisions in  $2^{69}$

Later (w/ coauthors): in  $2^{63}$

SHA-1 = **broken**

(phased out of SSL by 2017)

## Study Guide



### Definitions:

RandInt, Bernoulli  
experiment  
sample space, outcome  
event, probability  
conditioning  
Law of Total Probability  
independence

### Solving problems:

how to find probabilities  
how to condition  
proving independence