

# 15-251: Great Theoretical Ideas In Computer Science

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## Recitation 2

### Regular Announcements

- Homework Solution Sessions - Friday 5p-6p, Saturday 2p-3p in GHC 4301
- Small Groups this weekend
- Come to us if you had difficulties on Homework 1

### Definitions For All

- **Deterministic Finite Automaton (DFA):** A DFA  $M$  is a machine that reads a finite input one character at a time in one pass, transition from state to state, and ultimately accepts or rejects. Formally,  $M$  is a 5-tuple  $M = (Q, \Sigma, \delta, q_0, F)$ , where
  - $Q$  is a **finite, non-empty** set of states
  - $\Sigma$  is the **finite, non-empty** alphabet
  - $\delta : Q \times \Sigma \rightarrow Q$  is the transition function
  - $q_0 \in Q$  is the starting state
  - $F \subseteq Q$  is the set of accepting states
- **Regular language:** A language  $L$  is regular if  $L = L(M)$  for some DFA  $M$  ( $M$  decides  $L$ ).
- We have shown that if  $L_1$  and  $L_2$  are both regular languages over  $\Sigma^*$ , for some fixed  $\Sigma$ , then the following are all regular.
  - $\overline{L_1}$
  - $L_1 \cup L_2$
  - $L_1 \cap L_2$
  - $L_1 L_2$  (the concatenation of two regular languages)

### Odd Ones Out

Draw a DFA that decides the language

$$L = \{x : x \text{ has an even number of 1s and an odd number of 0s}\}$$

over the alphabet  $\Sigma = \{0, 1\}$ .

### #AllLanguagesAreBeautiful

Show that the following languages are irregular (but still beautiful).

- $L = \{a^{2^n} : n \in \mathbb{N}\}$
- $L = \{www : w \in \Sigma^*\}$

## States For Days

Define  $\mathcal{R}_n = \{x \mid x \in \{0,1\}^* \text{ and the } n^{\text{th}}\text{-symbol from the right is a } 1\}$ . Show that any DFA that accepts  $\mathcal{R}_n$  has at least  $2^n$  states.

## Double Trouble

Given a regular language  $L$  over some alphabet  $\Sigma$ , we define

$$K = \{x \mid xx \in L\}.$$

Prove that  $K$  is also regular.