## **Recitation 6**

#### Announcements

- Midterm 1 next Wednesday, February 28! It will be held in DH 2210 from 6.30pm to 9.30pm in place of the writing session. (Note the later end time.)
- We will be holding topical reviews with the venues and times to be confirmed. Watch Piazza for updates!

### Recap of some definitions and facts

- Stable Matching Problem
- Gale-Shapley Algorithm
- Regular graphs
- The tree-nity (the three salient features of a tree)
- Hamiltonian cycle
- The handshake lemma

### **Soulmates**

Call a man m and a woman w "soulmates" if they are paired with each other in every stable matching.

- (a) Given a man m and a woman w, design a polynomial-time algorithm to determine if they are soulmates.
- (b) Give a polynomial time algorithm to determine if an instance of the stable matching problem has a *unique* stable matching.

# "Clearly" Correct

A connected graph with no cycles is a tree. Consider the following claim and its proof.

**Claim:** Any graph with n vertices and n-1 edges is a tree.

**Proof:** We prove the claim by induction. The claim is clearly true for n = 1 and n = 2. Now suppose the claim holds for n = k. We'll prove that it also holds for n = k + 1. Let G be a graph with k vertices and k - 1 edges. By the induction hypothesis, G is a tree (and therefore clearly connected). Add a new vertex v to G by connecting it with any other vertex in G. So we create a new graph G' with k + 1 vertices and k edges. The new vertex we added is clearly a leaf, so it clearly does not create a cycle. Also, since G was connected, G' is clearly also connected. A connected graph with no cycles is a tree, so G' is also a tree. So the claim follows by induction.

Explain why the given proof is incorrect.

# 2 3 Proofs 4 You

Give two proofs (one using induction and another using a degree counting argument) for the following claim: the number of leaves in a tree with  $n \ge 2$  vertices is

$$2 + \sum_{\substack{v \in V \\ \deg(v) \ge 3}} (\deg(v) - 2)$$