## Theoretical Foundations of Computer Science



August 9, 2017

#### **Motivational Quote of the Day**

#### "Computer Science is no more about computers than astronomy is about telescopes."





#### PART I: What is (theoretical) computer science? How was it born?

#### PART 2: Uncomputable problems

#### PART I: What is computer science? How was it born?

#### What is **computer science**?

#### Is it:

"Writing (Python) programs that do certain tasks."

What is **theoretical computer science**?

# What is computer science?

Is it branch of:

- science?
- engineering?
- math?
- philosophy?
- sports?



# **Physics**

#### **Theoretical physics**

- come up with mathematical models Nature's language is mathematics
- derive the logical consequences

#### **Experimental physics**

- make observations about the universe
- test mathematical models with experiments

## **Applications/Engineering**



## The role of theoretical physics



#### **Theoretical Physics**

- science?
- engineering?
- math?
- philosophy?
- sports?



- **Computation**: manipulation of information/data.
- Algorithm: rigorous description of how the data is manipulated.

- **Computation**: manipulation of information/data.
- Algorithm: rigorous description of how the data is manipulated.

Input 
$$\rightarrow$$
 Calculator  $\rightarrow$  Output

- **Computation**: manipulation of information/data.
- Algorithm: rigorous description of how the data is manipulated.

- **Computation**: manipulation of information/data.
- Algorithm: rigorous description of how the data is manipulated.

- **Computation**: manipulation of information/data.
- Algorithm: rigorous description of how the data is manipulated.

Input 
$$\rightarrow$$
 **Evolution**  $\rightarrow$  Output

#### The computational lens



Computational physics Computational biology **Computational chemistry** Computational neuroscience **Computational economics** Computational finance **Computational linguistics** Computational statistics

# The role of theoretical computer science

Build a mathematical model for computation.

Explore the logical consequences. Gain insight about computation.

Look for interesting applications.



CMU undergrad



CMU Prof.



OK, we don't have everybody

http://youtu.be/pTeZP-XfuKI https://goo.gl/gGkpMv http://youtu.be/J4TkHuTmHsg

#### The role of theoretical computer science



## Simple examples of computation



Doing computation by following a simple algorithm.

## Simple examples of computation

Euclid's algorithm (~ 300BC):

```
def gcd(a, b):
    while (b != 0):
    t = b
        b = a % b
        a = t
```

return a

We have been using algorithms for thousands of years.

## Formalizing computation

We have been using algorithms for thousands of years.

Algorithm/Computation was only formalized in the 20th century!

Someone had to ask the right **question**.

## David Hilbert, 1900



#### **The Problems of Mathematics**

"Who among us would not be happy to lift the veil behind which is hidden the future; to gaze at the coming developments of our science and at the secrets of its development in the centuries to come? What will be the ends toward which the spirit of future generations of mathematicians will tend? What methods, what new facts will the new century reveal in the vast and rich field of mathematical thought?"

### Hilbert's 10th problem (1900)

Is there a finitary procedure to determine if a given multivariate polynomial with integral coefficients has an integral solution?

e.g. 
$$5x^2yz^3 + 2xy + y - 99xyz^4 = 0$$

## Entscheidungsproblem (1928)

Is there a finitary procedure to determine the validity of a given logical expression?

e.g. 
$$\neg \exists x, y, z, n \in \mathbb{N} : (n \ge 3) \land (x^n + y^n = z^n)$$

(Mechanization of mathematics)

#### Fortunately, the answer turned out to be NO.

# Gödel (1934):

Discusses some ideas for mathematical definitions of computation. But not confident what is a good definition.

# Church (1936):

Invents lambda calculus. Claims it should be the definition of an "algorithm".

# Gödel, Post (1936):

Arguments that Church's claim is not justified.

Meanwhile... in New Jersey... a certain British grad student, unaware of all these debates...







## Alan Turing (1936, age 22):

Describes a new model for computation, now known as the Turing Machine.™



#### Gödel, Kleene, and even Church:

"Umm. Yeah. He nailed it. Game over. "Algorithm" defined."

## Turing (1937): TMs ≡ lambda calculus

# Formalization of computation: Turing Machine

#### **Turing Machine:**



# **Church-Turing Thesis**

#### **Church-Turing Thesis:**

The intuitive notion of "computable" is captured by functions computable by a Turing Machine.

## (Physical) Church-Turing Thesis

Any computational problem that can be solved by a physical device, can be solved by a Turing Machine.



## Back to 2 of Hilbert's Problems

## Hilbert's 10th problem (1900)

Is there an **algorithm** (a TM) to determine if a given multivariate polynomial with integral coefficients has an integral solution?

e.g. 
$$5x^2yz^3 + 2xy + y - 99xyz^4 = 0$$

# Entscheidungsproblem (1928)

Is there an **algorithm** (a TM) to determine the validity of a given logical expression?

e.g. 
$$\neg \exists x, y, z, n \in \mathbb{N} : (n \ge 3) \land (x^n + y^n = z^n)$$

(Mechanization of mathematics)

## Back to 2 of Hilbert's Problems

## Hilbert's 10th problem (1900)

#### Matiyasevich-Robinson-Davis-Putnam (1970):



There is no algorithm to solve this problem.

## Entscheidungsproblem (1928)



# Turing (1936):

There is no algorithm to solve this problem.

- science?
- engineering?
- math?
- philosophy?
- sports?



# 2 Main Questions in TCS

#### **Computability** of a problem:

Is there an algorithm to solve it?

**Complexity** of a problem:

Is there an **efficient** algorithm to solve it?

- time

- space (memory)
- randomness
- quantum resources

# **Computational Complexity**

**Complexity** of a problem:

Is there an **efficient** algorithm to solve it?

- time
- space (memory)
- randomness
- quantum resources

#### 2 camps:

- trying to come up with efficient algorithms (algorithm designers)
- trying to show no efficient algorithm exists (complexity theorists)

# **Computational Complexity**

#### 2 camps:

- trying to come up with efficient algorithms (algorithm designers)
- trying to show no efficient algorithm exists (complexity theorists)

multiplying two integers factoring integers sorting a list protein structure prediction simulation of quantum systems computing Nash Equilibria of games

#### PART 2: Uncomputable problems



student submission

isPrime

the correct program isPrime

Do they return True on exactly the same inputs?





#### A "simpler" problem

Write an "autograder" that checks if a given program goes into an infinite loop.



## **Halting Problem**

**Inputs**: A Python program source code. An *input* to the program. ×

**Output**: True if the program halts for the given *input*. False otherwise.



**Theorem:** The **halting problem** is uncomputable.

Assume, for the sake of contradiction, such a program exists:

# def halt(program, inputToProgram): # program and inputToProgram are both strings



Assume, for the sake of contradiction, such a program exists:

def halt(program, inputToProgram):
 # program and inputToProgram are both strings

def turing(program):
 if (halt(program, program)):
 while True:
 pass # i.e. do nothing
 return False



Assume, for the sake of contradiction, such a program exists:

def halt(program, inputToProgram):
 # program and inputToProgram are both strings

def turing(program):
 if (halt(program, program)):
 while True:
 pass # i.e. do nothing
 return False

What happens when you call turing(turing)?



Assume, for the sake of contradiction, such a program exists:

def halt(program, inputToProgram):
 # program and inputToProgram are both strings

def turing(program):
 if (halt(program, program)):
 while True:
 pass # i.e. do nothing
 return False

What happens when you call turing(turing) ? if halt(turing, turing) is True ----> turing(turing) doesn't halt if halt(turing, turing) is False ----> turing(turing) halts

Assume, for the sake of contradiction, such a program exists:

def halt(program, inputToProgram):
 # program and inputToProgram are both strings

def turing(program):
 if (halt(program, program)):
 while True:
 pass # i.e. do nothing
 return False

What happens when you call turing(turing) ?

- No guaranteed autograder program. (;;)
- Consider the following program: **def fermat**():

t = 3while (True): **for** n **in** range(3, t+1): for x in range(1, t+1): for y in range(1, t+1): for z in range(1, t+1): **if**  $(x^{**}n + y^{**}n = z^{**}n)$ : **return** (x, y, z, n)t += 1

**Question**: Does this program halt?



- Consider the following program (written in MAPLE):

```
numberToTest := 2;
flag := 1;
while flag = 1 do
  flag := 0;
  numberToTest := numberToTest + 2;
  for p from 2 to numberToTest do
    if isPrime(p) and isPrime(numberToTest-p) then
       flag := 1;
       break; #exits the for loop
                                               Goldbach
    end if
                                              Conjecture
  end for
end do
```

#### **Question**: Does this program halt?

- Reductions to other problems imply that those problems are uncomputable as well.

#### Entscheidungsproblem

Is there an algorithm to determine the validity of a given logical expression?

e.g.  $\neg \exists x, y, z, n \in \mathbb{N} : (n \ge 3) \land (x^n + y^n = z^n)$ 

#### **Hilbert's 10th Problem**

Is there an algorithm to determine if a given multivariate polynomial with integral coefficients has an integral solution?

# **Input**: A finite collection of "dominoes" having strings written on each half.



**Output**: **True** if it is possible to match the strings.



#### **Uncomputable!**

Proved in 1946 by Post.

**Input**: A finite collection of "Wang Tiles" (squares) with colors on the edges.



Output: True if it is possible to make an infinite grid from copies of the given squares, where touching sides must color-match.

**Uncomputable!** 

Proved in 1966 by Berger.

**Input**: Two 21x21 matrices of integers A and B.

**Output**: **True** iff it is possible to multiply A and B together (multiple times in any order) to get to the 0 matrix.

#### **Uncomputable!**

Proved in 2007 by Halava, Harju, Hirvensalo.

Different laws of physics ---->

Different computational devices ---->

Every problem computable (???)

Can you come up with sensible laws of physics such that the Halting Problem becomes computable?

That was about the basic question on whether every problem is computable.

# Some other interesting questions in TCS

#### **Time vs Space**

If a problem has a space-efficient solution does it also have a time-efficient solution?

#### **Power of randomness**

Can every randomized algorithm be derandomized efficiently?

#### Power of quantum information

Can we use quantum properties of matter to build faster computers?

#### P vs NP